IMPROVED SMOOTHING-TYPE METHODS FOR THE SOLUTION OF LINEAR PROGRAMS¹

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Abstract. We consider a smoothing-type method for the solution of linear programs. Its main idea is to reformulate the primal-dual optimality conditions as a nonlinear and nonsmooth system of equations, and to apply a Newton-type method to a smooth approximation of this nonsmooth system. The method presented here is a predictor-corrector method, and is closely related to some methods recently proposed by Burke and Xu on the one hand, and by the authors on the other hand. However, here we state stronger global and/or local convergence properties. Moreover, we present quite promising numerical results for the whole netlib test problem collection.

Key Words. Linear programs, smoothing, predictor-corrector method, Newton's method, global convergence, quadratic convergence.

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1 Introduction

In this paper we describe an algorithm for the solution of linear programs given either in primal form

$$\min c^T x \quad \text{s.t.} \quad Ax = b, \ x \ge 0 \tag{1}$$

or in dual form

$$\max b^T \lambda \quad \text{s.t.} \quad A^T \lambda + s = c, \ s \ge 0, \tag{2}$$

where $A \in \mathbb{R}^{m \times n}$, $c \in \mathbb{R}^n$, and $b \in \mathbb{R}^m$ are the given data, and A is assumed to have full rank throughout this paper. The idea of our algorithm is to solve the corresponding optimality conditions

$$A^{T}\lambda + s = c,$$

 $Ax = b,$
 $x_{i} \ge 0, s_{i} \ge 0, x_{i}s_{i} = 0 \quad \forall i = 1, ..., n.$ (3)

To this end, note that (3) is totally equivalent to (1) and (2) in the sense that (3) has a solution if and only if (1) or (2) has a solution.

The widely used class of primal-dual interior-point methods follow a similar idea: They are also based on the optimality conditions (3) and introduce a certain perturbation of (3) depending on a parameter $\tau > 0$:

$$A^{T}\lambda + s = c,$$

 $Ax = b,$
 $x_{i} > 0, s_{i} > 0, x_{i}s_{i} = \tau^{2} \quad \forall i = 1, ..., n.$

$$(4)$$

The system (4) is usually called the *central path conditions*, and is parameterized here by τ^2 instead of τ just for technical reasons which will become clear in Section 2.

Under certain assumptions, there is a unique solution $w_{\tau} = (x_{\tau}, \lambda_{\tau}, s_{\tau})$ of (4) for each $\tau > 0$. The corresponding mapping

$$\tau \mapsto w_{\tau}$$

is called the *central path*, and the main idea of interior-point methods is to follow this central path numerically. This is typically done by applying Newton's method to the *equations* within the central path conditions (4), whereas a suitable stepsize rule takes care of the strict inequality constraints. In particular, all iterates generated by a primal-dual interior-point method satisfy the strict inequality constraints.

Smoothing-type methods follow a different approach. The general idea of these methods is to reformulate the optimality conditions (3) as a system of equations (not involving any inequalities). Since this system is typically nonsmooth, it then gets approximated by a smooth system of equations to which Newton's method can be applied, see [3, 5, 11, 16] and references therein for a couple of examples following this pattern.

The method to be presented here follows an idea by Jiang [12] and is based on a smooth equation reformulation of the optimality conditions (3) themselves. It is, however, closely related to both smoothing-type methods and interior-point methods. This will be made clear as soon as we develop the algorithm in Section 2.

The algorithm we present in this paper is quite similar to the predictor-corrector method recently proposed by Burke and Xu [2], see also [1]. Our method is also a predictor-corrector

method, with the corrector step being responsible for the global convergence and the predictor step guaranteeing local fast convergence under suitable assumptions. In fact, our corrector step is identical to the one by Burke and Xu [2], but we prove a different global convergence result for it using less stringent assumptions. On the other hand, the predictor step we use here is taken from [9] (and was essentially introduced by Chen, Qi and Sun [6]) and can be shown to be locally quadratically convergent under weaker assumptions than those used by Burke and Xu [2].

We therefore view our algorithm as an improved smoothing-type method due to some better theoretical properties if compared with the method by Burke and Xu [2]. In addition, it also improves the authors' previous method [9] due to some stronger global convergence properties. Furthermore, the numerical results seem to be somewhat better than for the corresponding method in [9].

This paper is organized as follows: In Section 2, we develop our algorithm and give a detailed statement. The global and local convergence properties are investigated in Section 3. Extensive numerical results are presented in Section 4, and Section 5 concludes this paper with some final remarks.

The notation used in this paper is rather standard: \mathbb{R}^n denotes the n-dimensional real vector space. For $x \in \mathbb{R}^n$, we use the subscript x_i in order to indicate the ith component of x, whereas a superscript like in x^k is used to indicate that this is the kth iterate of a sequence $\{x^k\} \subseteq \mathbb{R}^n$. Quite often, we will consider a triple of the form $w = (x^T, \lambda^T, s^T)^T$, where $x \in \mathbb{R}^n$, $\lambda \in \mathbb{R}^m$, and $s \in \mathbb{R}^n$; of course, w is then a vector in \mathbb{R}^{n+m+n} . In order to simplify our notation, however, we will usually write $w = (x, \lambda, s)$ instead of using the mathematically more correct formula $w = (x^T, \lambda^T, s^T)^T$. If $x, y \in \mathbb{R}^n$ are any given vectors satisfying the inequality $x_i \geq y_i$ for all indices $i = 1, \ldots, n$, we simply write $x \geq y$. Finally, the symbol $\|\cdot\|$ is used for the Euclidean vector norm, whereas $\|\cdot\|_{\infty}$ indicates the maximum norm.

2 Development of Algorithm

This section is devoted to the development of our algorithm. To this end, let $\varphi : \mathbb{R}^2 \to \mathbb{R}$ always denote the minimum function

$$\varphi(a,b) := 2\min\{a,b\},\$$

with the factor 2 being used here only for cosmetical reasons. Define

$$\Phi(w) := \Phi(x, \lambda, s) := \begin{pmatrix} A^T \lambda + s - c \\ Ax - b \\ \phi(x, s) \end{pmatrix},$$

where

$$\phi(x,s) := (\varphi(x_1,s_1),\ldots,\varphi(x_n,s_n))^T \in \mathbb{R}^n.$$

Since φ is an NCP-function, i.e.,

$$\varphi(a,b) = 0 \iff a \ge 0, b \ge 0, ab = 0,$$

the following observation is an immediate consequence of the definition of the mapping Φ :

$$w^* = (x^*, \lambda^*, s^*)$$
 solves (3) $\iff w^*$ solves $\Phi(w) = 0$.

However, the system $\Phi(w)=0$ is nonsmooth. Therefore, let $\varphi_{\tau}:\mathbb{R}^2\to\mathbb{R}$ denote the smoothed minimum function

$$\varphi_{\tau}(a,b) := a + b - \sqrt{(a-b)^2 + 4\tau^2},$$

typically called the Chen-Harker-Kanzow-Smale smoothing function [4, 13, 15], where $\tau > 0$ is the smoothing parameter. Then define

$$\Phi_{\tau}(w) := \Phi_{\tau}(x, \lambda, s) := \begin{pmatrix} A^{T}\lambda + s - c \\ Ax - b \\ \phi_{\tau}(x, s) \end{pmatrix},$$

where

$$\phi_{\tau}(x,s) := (\varphi_{\tau}(x_1,s_1), \dots, \varphi_{\tau}(x_n,s_n))^T \in \mathbb{R}^n.$$

It was observed in [13] that the following equivalence holds:

$$w_{\tau} = (x_{\tau}, \lambda_{\tau}, s_{\tau}) \text{ solves } (4) \iff w_{\tau} \text{ solves } \Phi_{\tau}(w) = 0,$$

i.e., we obtain a reformulation of the central path conditions as a nonlinear and smooth system of equations in this way.

So far, we have viewed τ as a parameter. In the sequel, however, it will sometimes be useful to view τ as an independent variable. In order to make this different point of view clear in our notation, let us write

$$\theta(x, s, \tau) := \phi_{\tau}(x, s).$$

Moreover, we will exploit the mapping $\Theta : \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^n \times [0, \infty) \to \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^n \times \mathbb{R}$ defined by

$$\Theta(w,\tau) := \Theta(x,\lambda,s,\tau) := \begin{pmatrix} A^T \lambda + s - c \\ Ax - b \\ \theta(x,s,\tau) \\ \tau \end{pmatrix}.$$

Note that the definition of $\Theta(w,\tau)$ is not equal to $\Phi_{\tau}(w)$ since we have added one more line. Since the equation $\Theta(w,\tau)=0$ automatically implies $\tau=0$, we obtain the equivalence

$$w^* = (x^*, \lambda^*, s^*)$$
 solves (3) \iff $(w^*, 0)$ solves $\Theta(w, \tau) = 0$.

In this way we therefore get a reformulation of the optimality conditions (3) with τ being viewed as an independent variable. This kind of reformulation goes back to Jiang [12].

For later reference, it will be important to exploit the relation between Newton's method applied to the system $\Phi_{\tau}(w) = 0$ and applied to the system $\Theta(w, \tau) = 0$. First consider the

system $\Phi_{\tau}(w) = 0$, and assume that $w^k = (x^k, \lambda^k, s^k)$ denotes the current iterate and $\tau_k > 0$ the current value of the smoothing parameter. Then we define

$$w^{k+1} = w^k + t_k \Delta w^k$$

for a suitable stepsize $t_k > 0$, where the correction vector $\Delta w^k = (\Delta x^k, \Delta \lambda^k, \Delta s^k)$ is a solution of the linear system of equations

$$\Phi_{\tau_k}'(w^k)\Delta w^k = -\Phi_{\tau_k}(w^k).$$

Taking into account the definition of Φ_{τ} , this equation is equivalent to

$$\begin{pmatrix} 0 & A^T & I \\ A & 0 & 0 \\ D_{a,\tau}^k & 0 & D_{b,\tau}^k \end{pmatrix} \begin{pmatrix} \Delta x \\ \Delta \lambda \\ \Delta s \end{pmatrix} = \begin{pmatrix} -A^T \lambda^k - s^k + c \\ -Ax^k + b \\ -\phi_{\tau_k}(x^k, s^k) \end{pmatrix}, \tag{5}$$

where

$$D_{a,\tau}^k := \operatorname{diag}\left(\frac{\partial \varphi_{\tau_k}}{\partial a}(x_1^k, s_1^k), \dots, \frac{\partial \varphi_{\tau_k}}{\partial a}(x_n^k, s_n^k)\right) \in \mathbb{R}^{n \times n}$$

and, similarly,

$$D_{b,\tau}^k := \operatorname{diag}\left(\frac{\partial \varphi_{\tau_k}}{\partial b}(x_1^k, s_1^k), \dots, \frac{\partial \varphi_{\tau_k}}{\partial b}(x_n^k, s_n^k)\right) \in \mathbb{R}^{n \times n}.$$

On the other hand, if we apply Newton's method to the system $\Theta(w,\tau)=0$, we have to solve an equation like

$$\Theta'(w^k, \tau_k) \left(\begin{array}{c} \Delta w \\ \Delta \tau \end{array}\right) = -\Theta(w^k, \tau_k)$$

at each iteration, where the derivatives are taken with respect to w and τ . Hence the above system becomes

$$\begin{pmatrix} 0 & A^{T} & I & 0 \\ A & 0 & 0 & 0 \\ D_{a,\tau}^{k} & 0 & D_{b,\tau}^{k} & d_{\tau}^{k} \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} \Delta x \\ \Delta \lambda \\ \Delta s \\ \Delta \tau \end{pmatrix} = \begin{pmatrix} -A^{T}\lambda^{k} - s^{k} + c \\ -Ax^{k} + b \\ -\theta(x^{k}, s^{k}, \tau_{k}) \\ -\tau_{k} \end{pmatrix},$$
(6)

where

$$d_{\tau}^{k} := \left(\frac{\partial \theta}{\partial \tau}(x_{1}^{k}, s_{1}^{k}, \tau_{k}), \dots, \frac{\partial \theta}{\partial \tau}(x_{n}^{k}, s_{n}^{k}, \tau_{k})\right)^{T} \in \mathbb{R}^{n}.$$

Motivated by similar considerations in the field of interior-point methods (see, e.g., Wright [18]), we will consider a generalization of the system (6) and replace the parameter τ_k on the last line of the right-hand side in (6) by $\sigma_k \tau_k$ for some number $\sigma_k \in [0, 1]$, i.e., we solve

$$\begin{pmatrix} 0 & A^{T} & I & 0 \\ A & 0 & 0 & 0 \\ D_{a,\tau}^{k} & 0 & D_{b,\tau}^{k} & d_{\tau}^{k} \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} \Delta x \\ \Delta \lambda \\ \Delta s \\ \Delta \tau \end{pmatrix} = \begin{pmatrix} -A^{T}\lambda^{k} - s^{k} + c \\ -Ax^{k} + b \\ -\theta(x^{k}, s^{k}, \tau_{k}) \\ -\sigma_{k}\tau_{k} \end{pmatrix}$$
(7)

(the choice $\sigma_k = 1$ corresponds to (6)). Note, however, that we do not replace τ_k by $\sigma_k \tau_k$ in the definition of the function $\theta(x, s, \tau)$. In order to have a short-hand notation for the linear system (7), we introduce the function

$$\Theta_{\sigma}(w,\tau) := \begin{pmatrix} A^{T}\lambda + s - c \\ Ax - b \\ \theta(x,s,\tau) \\ \sigma\tau \end{pmatrix}$$

with the subscript σ indicating the dependence of Θ on the parameter σ . Then the linear system (7) can be rewritten as

$$\Theta'(w^k, \tau_k) \begin{pmatrix} \Delta w \\ \Delta \tau \end{pmatrix} = -\Theta_{\sigma}(w^k, \tau_k).$$

Note that this or, equivalently, (7) immediately gives

$$\Delta \tau_k = -\sigma_k \tau_k. \tag{8}$$

Replacing this expression into the remaining equations of (7), we obtain

$$\begin{pmatrix} 0 & A^T & I \\ A & 0 & 0 \\ D_{a,\tau}^k & 0 & D_{b,\tau}^k \end{pmatrix} \begin{pmatrix} \Delta x \\ \Delta \lambda \\ \Delta s \end{pmatrix} = \begin{pmatrix} -A^T \lambda^k - s^k + c \\ -Ax^k + b \\ -\theta(x^k, s^k, \tau_k) + \sigma_k \tau_k d_\tau^k \end{pmatrix}. \tag{9}$$

Obviously, this can be rewritten as

$$\Phi_{\tau_k}'(w^k)\Delta w = -\Phi_{\tau_k}(w^k) + \sigma_k \tau_k \begin{pmatrix} 0 \\ 0 \\ d_{\tau}^k \end{pmatrix}.$$

This shows that the linear system (7) can be viewed as a perturbation of the system (5), with the perturbation being active only in the third block row of the right-hand side.

Based on the notation introduced so far, we next give a precise statement of our predictorcorrector smoothing method.

Algorithm 2.1 (Predictor-Corrector Smoothing Method)

- (S.0) (Initialization) Choose $w^0 := (x^0, \lambda^0, s^0) \in \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^n$ such that $A^T \lambda^0 + s^0 = c$, $Ax^0 = b$, choose $\tau_0 > 0$, select $\beta \ge \|\Phi_{\tau_0}(w^0)\|/\tau_0$, $\rho \in (0, 1)$, $0 < \hat{\sigma}_{\min} < \hat{\sigma}_{\max} < 1$, $\varepsilon \ge 0$, and set k := 0.
- (S.1) (Termination Criterion) If $\|\Phi(w^k)\| < \varepsilon$: STOP.
- (S.2) (Predictor Step) Compute a solution $(\Delta w^k, \Delta \tau_k) = (\Delta x^k, \Delta \lambda^k, \Delta s^k, \Delta \tau_k) \in \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}$ of the linear system

$$\Theta'(w^k, \tau_k) \begin{pmatrix} \Delta w \\ \Delta \tau \end{pmatrix} = -\Theta(w^k, 0). \tag{10}$$

If $\|\Phi(w^k + \Delta w^k)\| = 0$: STOP. Otherwise, if

$$\|\theta(x^k + \Delta x^k, s^k + \Delta s^k, \tau_k)\| > \beta \tau_k,$$

then set

$$\hat{w}^k := w^k, \ \hat{\tau}_k := \tau_k, \ \eta_k := 1,$$

else compute $\eta_k = \rho^{\ell_k}$, where ℓ_k is the nonnegative integer such that

$$\|\theta(x^k + \Delta x^k, s^k + \Delta s^k, \rho^j \tau_k)\| \le \beta \rho^j \tau_k \quad \forall j = 0, 1, 2, \dots, \ell_k \text{ and } \|\theta(x^k + \Delta x^k, s^k + \Delta s^k, \rho^{\ell_k + 1} \tau_k)\| > \beta \rho^{\ell_k + 1} \tau_k,$$

and set $\hat{\tau}_k := \eta_k \tau_k$ and

$$\hat{w}^k := \begin{cases} w^k & \text{if } \ell_k = 0, \\ w^k + \Delta w^k & \text{otherwise.} \end{cases}$$

(S.3) (Corrector Step)

Choose $\hat{\sigma}_k \in [\hat{\sigma}_{\min}, \hat{\sigma}_{\max}]$, and compute a solution $(\Delta \hat{w}^k, \Delta \hat{\tau}_k) = (\Delta \hat{x}^k, \Delta \hat{\lambda}^k, \Delta \hat{s}^k, \Delta \hat{\tau}_k) \in \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}$ of the linear system

$$\Theta'(\hat{w}^k, \hat{\tau}_k) \begin{pmatrix} \Delta \hat{w} \\ \Delta \hat{\tau} \end{pmatrix} = -\Theta_{\hat{\sigma}_k}(\hat{w}^k, \hat{\tau}_k). \tag{11}$$

Let $\hat{t}_k = \max\{\rho^{\ell} | \ell = 0, 1, 2, ...\}$ such that

$$\|\theta(\hat{x}^k + \hat{t}_k \Delta \hat{x}^k, \hat{s}^k + \hat{t}_k \Delta \hat{s}^k, (1 - \hat{\sigma}_k \hat{t}_k) \hat{\tau}_k)\| \le \beta (1 - \hat{\sigma}_k \hat{t}_k) \hat{\tau}_k. \tag{12}$$

Set $w^{k+1} := \hat{w}^k + \hat{t}_k \Delta \hat{w}^k$ and $\tau_{k+1} := (1 - \hat{\sigma}_k \hat{t}_k) \hat{\tau}_k$.

(S.4) (Update) Set $k \leftarrow k+1$, and go to Step (S.1).

To get a better understanding of the way Algorithm 2.1 works, let us add a couple of comments. In Step (S.0), we require the starting point $w^0 = (x^0, \lambda^0, s^0)$ to be feasible with respect to the linear equations $A^T \lambda + s = c$ and Ax = b. Since the components x^0 and s^0 do not have to be positive (like in interior-point methods), it is relatively easy to find such a starting point.

In the predictor step, we first compute a search direction by solving the linear system (10). The interesting part about this linear system is the fact that the right-hand side of (10) is unperturbed with respect to τ , whereas we use the standard Jacobian of the perturbed function $\Theta(w,\tau)$ on the left. This may be viewed as the counterpart of the affine scaling step typically used as a predictor in primal-dual interior-point methods, see Wright [18]. Like in the interior-point setting, this predictor step will eventually guarantee local fast convergence (under suitable assumptions).

After having computed the search direction in (10), we try to reduce the smoothing parameter τ_k as much as possible with the only restriction that the full step stays within a certain neighbourhood of the central path, cf. Lemma 3.2 (c) below.

While our predictor step is different from the one used by Burke and Xu [2] in their smoothing-type method, the corrector step in (S.3) coincides with the one from [2]. Note that the linear system (11) is precisely the one from (7) and includes a perturbation on the right-hand side as well. The predictor step also contains a procedure to reduce the smoothing parameter. This procedure will guarantee global convergence in the sense that τ_k decreases to zero under mild conditions.

3 Convergence Properties

This section investigates the global and local convergence properties of Algorithm 2.1. To this end, we assume throughout this section that the termination parameter ε is equal to zero, and that Algorithm 2.1 generates an infinite number of iterates w^k , i.e., we assume that we do not stop after a finite number of iterations in a point w^k satisfying the optimality conditions (3).

We first note that Algorithm 2.1 is well-defined.

Lemma 3.1 The following statements hold for any $k \in \mathbb{N}$:

- (a) The linear systems (10) and (11) have a unique solution.
- (b) There is a unique η_k satisfying the conditions in Step (S.2).
- (c) The stepsize \hat{t}_k in (S.3) is uniquely defined.

Consequently, Algorithm 2.1 is well-defined.

Proof. (a) The structure of the Jacobian $\Theta'(w^k, \tau_k)$ in (6) shows that this matrix is nonsingular if and only if $\Phi'_{\tau_k}(w^k)$ is nonsingular. The latter, however, was noted in [8, Proposition 3.1].

- (b) This statement follows from [9, Proposition 3.2] and is essentially due to Burke and Xu [2].
- (c) This result follows from the proof of Theorem 1 in Burke and Xu [2]. \Box

We next state some simple properties of Algorithm 2.1 to which we will refer a couple of times in our subsequent analysis.

Lemma 3.2 The sequences $\{w^k\} = \{(x^k, \lambda^k, s^k)\}$ and $\{\tau_k\}$ generated by Algorithm 2.1 have the following properties:

- (a) $A^{T}\lambda^{k} + s^{k} = c$ and $Ax^{k} = b$ for all $k \in \mathbb{N}$.
- (b) $\tau_k = \tau_0(1 \hat{\sigma}_0\hat{t}_0)\eta_0\cdots(1 \hat{\sigma}_{k-1}\hat{t}_{k-1})\eta_{k-1} \text{ for all } k \in \mathbb{N}.$
- (c) $\|\theta(x^k, s^k, \tau_k)\| \le \beta \tau_k$ for all $k \in \mathbb{N}$.

Proof. (a) For k = 0, this follows from the choice of our starting point in Step (S.0). Newton's method then guarantees that the linear equations $A^T \lambda + s = c$ and Ax = b are also satisfied for all $k \geq 1$.

(b) Step (S.2) of Algorithm 2.1 implies that $\hat{\tau}_k = \eta_k \tau_k$. The updating rules in Step (S.3) therefore give

$$\tau_{k+1} = (1 - \hat{\sigma}_k \hat{t}_k) \hat{\tau}_k = (1 - \hat{\sigma}_k \hat{t}_k) \eta_k \tau_k$$

for all $k \in \mathbb{N}$. This gives the desired formula, see also [2, Theorem 2].

(c) The choice of the starting point $w^0 = (x^0, \lambda^0, s^0)$ and β in Step (S.0) guarantee that we have $\|\theta(x^k, s^k, \tau_k)\| \leq \beta \tau_k$ for k = 0. The updating rules in Step (S.3) show that this inequality holds for all $k \in \mathbb{N}$.

The next result is quite simple and will be used in order to show that the sequence $\{w^k\}$ generated by Algorithm 2.1 will be bounded under certain conditions.

Lemma 3.3 The sequences $\{w^k\} = \{(x^k, \lambda^k, s^k)\}$ and $\{\tau_k\}$ generated by Algorithm 2.1 satisfy the inequality

$$\|\min\{x^k, s^k\}\|_{\infty} \le \kappa \tau_k$$

for all $k \in \mathbb{N}$ with $\kappa := (2 + \beta)/2$.

Proof. Let θ_i denote the *i*th component function of θ , i.e.,

$$\theta_i(a, b, \tau) := a + b - \sqrt{(a - b)^2 + 4\tau^2}.$$

Then it is easy to see that the inequality

$$|\theta_i(a, b, 0) - \theta_i(a, b, \tau)| \le 2\tau$$

holds for all $a, b \in \mathbb{R}$ and all $\tau > 0$. Using Lemma 3.2 (c), it then follows that

$$2|\min\{x_{i}^{k}, s_{i}^{k}\}| = |\theta_{i}(x_{i}^{k}, s_{i}^{k}, 0)|$$

$$\leq |\theta_{i}(x_{i}^{k}, s_{i}^{k}, \tau_{k})| + |\theta_{i}(x_{i}^{k}, s_{i}^{k}, 0) - \theta_{i}(x_{i}^{k}, s_{i}^{k}, \tau_{k})|$$

$$\leq ||\theta(x^{k}, s^{k}, \tau_{k})|| + |\theta_{i}(x_{i}^{k}, s_{i}^{k}, 0) - \theta_{i}(x_{i}^{k}, s_{i}^{k}, \tau_{k})|$$

$$\leq (\beta + 2)\tau_{k}$$

for all $k \in \mathbb{N}$ and all i = 1, ..., n. This implies

$$\|\min\{x^k, s^k\}\|_{\infty} \le \kappa \tau_k$$

for all $k \in \mathbb{N}$, where κ denotes the constant specified in the statement of our lemma.

We next show that the sequence $\{w^k\}$ generated by Algorithm 2.1 remains bounded provided that there is a strictly feasible point for the optimality conditions (3) (i.e., a vector $\hat{w} = (\hat{x}, \hat{\lambda}, \hat{s})$ satisfying $A^T \hat{\lambda} + \hat{s} = c, A\hat{x} = b$ and $\hat{x} > 0, \hat{s} > 0$) and that the initial smoothing parameter $\tau_0 > 0$ is sufficiently small. This boundedness result is similar to one given by Chen and Ye [7] in the context of box constrained variational inequality problems.

Proposition 3.4 Assume that there is a strictly feasible point $(\hat{x}, \hat{\lambda}, \hat{s})$ for the optimality conditions (3), and suppose that the initial smoothing parameter $\tau_0 > 0$ satisfies

$$\tau_0 < \frac{1}{\kappa} \min_{i=1,\dots,n} \{\hat{x}_i, \hat{s}_i\},$$

where $\kappa := (2 + \beta)/2$ denotes the constant from Lemma 3.3. Then the sequence $\{w^k\} = \{(x^k, \lambda^k, s^k)\}$ generated by Algorithm 2.1 is bounded.

Proof. Assume that the sequence $\{w^k\} = \{(x^k, \lambda^k, s^k)\}$ generated by Algorithm 2.1 is unbounded. Since $\{\tau_k\}$ is monotonically decreasing, it follows from Lemma 3.3 that

$$|\min\{x_i^k, s_i^k\}| \le \|\min\{x^k, s^k\}\|_{\infty} \le \kappa \tau_k \le \kappa \tau_0$$
 (13)

for all $k \in \mathbb{N}$ and all i = 1, ..., n. This obviously implies that there is no index $i \in \{1, ..., n\}$ such that $x_i^k \to -\infty$ or $s_i^k \to -\infty$ on a subsequence. Therefore, all components of the two sequences $\{x^k\}$ and $\{s^k\}$ are bounded from below.

On the other hand, the sequence $\{w^k\} = \{(x^k, \lambda^k, s^k)\}$ is unbounded by assumption. This implies that there is at least one component $i \in \{1, \ldots, n\}$ such that $x_i^k \to +\infty$ or $s_i^k \to +\infty$ on a subsequence since otherwise the two sequences $\{x^k\}$ and $\{s^k\}$ would be bounded which, in turn, would imply the boundedness of the sequence $\{\lambda^k\}$ as well because we have $A^T\lambda^k + s^k = c$ for all $k \in \mathbb{N}$ (cf. Lemma 3.2 (a)) and because A is assumed to have full rank.

Now let $\hat{w} = (\hat{x}, \hat{\lambda}, \hat{s}) \in \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^n$ be the strictly feasible point from our assumption. Then, in particular, we have

$$A^T \hat{\lambda} + \hat{s} = c$$
 and $A\hat{x} = b$.

Since we also have

$$A^T \lambda^k + s^k = c$$
 and $Ax^k = b$

for all $k \in \mathbb{N}$ by Lemma 3.2 (a), we get

$$A^{T}(\hat{\lambda} - \lambda^{k}) + (\hat{s} - s^{k}) = 0$$
 and $A(\hat{x} - x^{k}) = 0$ (14)

by subtracting these equations. Premultiplying the first equation in (14) with $(\hat{x} - x^k)^T$ and taking into account the second equation in (14) gives

$$\sum_{i=1}^{n} (\hat{x}_i - x_i^k)(\hat{s}_i - s_i^k) = (\hat{x} - x^k)^T (\hat{s} - s^k) = 0.$$
 (15)

We now assume without loss of generality that there is at least one component i such that $\{x_i^k\}$ is unbounded, i.e., $\{x_i^k\}_K \to +\infty$ for a suitable subset $K \subseteq \mathbb{N}$ (the argument would be similar if there would exist at least one component i with $\{s_i^k\}$ being unbounded). Let us define the following index sets:

 $I_x := \{i \mid \{x_i^k\}_K \text{ is unbounded}\},\$ $I_s := \{i \mid \{s_i^k\}_K \text{ is unbounded}\},\$ $I_b := \{i \mid \{x_i^k\}_K \text{ and } \{s_i^k\}_K \text{ are bounded}\}.$

Note that I_x is nonempty, whereas I_s (and I_b) might be empty. Using the definitions of these three index sets and subsequencing if necessary, we obtain from (13) that

$$\{x_i^k\}_K \to +\infty \quad \text{and} \quad s_i^k \le \kappa \tau_0 \ \forall k \in K \ \forall i \in I_x$$
 (16)

and

$$\{s_i^k\}_K \to +\infty \quad \text{and} \quad x_i^k \le \kappa \tau_0 \ \forall k \in K \ \forall i \in I_s,$$
 (17)

whereas there is a constant $c \in \mathbb{R}$ such that

$$\sum_{i \in I_b} (x_i^k - \hat{x}_i)(s_i^k - \hat{s}_i) \le c$$

for all $k \in K$. Using (15) then gives

$$c \geq \sum_{i \in I_b} (x_i^k - \hat{x}_i)(s_i^k - \hat{s}_i)$$

$$= \sum_{i \in I_x} (x_i^k - \hat{x}_i)(\hat{s}_i - s_i^k) + \sum_{i \in I_s} (x_i^k - \hat{x}_i)(\hat{s}_i - s_i^k)$$

for all $k \in K$. However, the right-hand side is unbounded on a subsequence due to (16) and (17) since $\hat{s}_i - s_i^k \ge \hat{s}_i - \kappa \tau_0 > 0$ ($i \in I_x$) and $\hat{x}_i - x_i^k \ge \hat{x}_i - \kappa \tau_0 > 0$ ($i \in I_s$) in view of our choice of $\tau_0 > 0$. This contradiction completes the proof.

Note that Proposition 3.4 guarantees the boundedness of the iterates w^k provided that the initial smoothing parameter is sufficiently small. On the other hand, it is interesting to note that Burke and Xu [2] can prove the boundedness of their iterates under the assumption that τ_0 is sufficiently large. In fact, Burke and Xu [2] can provide a lower bound for their choice of τ_0 which is known a priori, whereas our upper bound from Proposition 3.4 is, in general, not known. However, the lower bound from [2] could be very large, and this, in turn, could have a bad influence on the numerical behaviour of the smoothing-type method. — In any case, it should be noted that some interior-point methods generate bounded iterates under the sole assumption that the primal and dual linear programs (1) and (2), respectively, are feasible (rather than strictly feasible).

We next give a global convergence result for Algorithm 2.1. Note that this result is different from the one provided by Burke and Xu [2]. (They use an assumption which is even stronger than the one we use for our local convergence result in Theorem 3.8; on the other hand, the main emphasis in [2] was to prove a global linear rate of convergence result.)

Theorem 3.5 Assume that the sequence $\{w^k\} = \{(x^k, \lambda^k, s^k)\}$ generated by Algorithm 2.1 has at least one accumulation point. Then $\{\tau_k\}$ converges to zero.

Proof. Since the sequence $\{\tau_k\}$ is monotonically decreasing and bounded from below by zero, it converges to a number $\tau_* \geq 0$. If $\tau_* = 0$, we are done.

So assume that $\tau_* > 0$. Then the updating rules in Step (S.2) of Algorithm 2.1 immediately give

$$\hat{w}^k = w^k, \quad \hat{\tau}_k = \tau_k, \quad \text{and} \quad \eta_k = 1$$
 (18)

for all $k \in \mathbb{N}$ sufficiently large. Subsequencing if necessary, we assume without loss of generality that (18) holds for all $k \in \mathbb{N}$. Then Lemma 3.2 (b) and $\hat{\sigma}_k \geq \hat{\sigma}_{\min}$ yield

$$\tau_k = \tau_0 \prod_{j=0}^{k-1} (1 - \hat{\sigma}_j \hat{t}_j) \le \tau_0 \prod_{j=0}^{k-1} (1 - \hat{\sigma}_{\min} \hat{t}_j). \tag{19}$$

Since $\tau_k \to \tau_* > 0$ by assumption, it follows from (19) that $\lim_{k\to\infty} \hat{t}_k = 0$. Therefore, the stepsize $\hat{\alpha}_k := \hat{t}_k/\rho$ does not satisfy the line search criterion (12) for all $k \in \mathbb{N}$ sufficiently large. Hence we have

$$\|\theta(\hat{x}^k + \hat{\alpha}_k \Delta \hat{x}^k, \hat{s}^k + \hat{\alpha}_k \Delta \hat{s}^k, (1 - \hat{\sigma}_k \hat{\alpha}_k) \hat{\tau}_k)\| > \beta (1 - \hat{\sigma}_k \hat{\alpha}_k) \hat{\tau}_k$$
 (20)

for all these $k \in \mathbb{N}$.

Now let $w^* = (x^*, \lambda^*, s^*)$ be an accumulation point of the sequence $\{w^k\}$, and let $\{w^k\}_K$ be a subsequence converging to w^* . Since $\hat{\sigma}_k \in [\hat{\sigma}_{\min}, \hat{\sigma}_{\max}]$ for all $k \in \mathbb{N}$, we can assume without loss of generality that the subsequence $\{\hat{\sigma}_k\}_K$ converges to some number $\hat{\sigma}_* \in [\hat{\sigma}_{\min}, \hat{\sigma}_{\max}]$. Furthermore, since $\tau_* > 0$, it follows from Lemma 3.1 (a) that the corresponding subsequence $\{(\Delta \hat{w}^k, \Delta \hat{\tau}_k)\}_K$ converges to a vector $(\Delta \hat{w}^*, \Delta \hat{\tau}_*) = (\Delta \hat{x}^*, \Delta \hat{\lambda}^*, \Delta \hat{s}^*, \Delta \hat{\tau}_*)$, where $(\Delta \hat{w}^*, \Delta \hat{\tau}_*)$ is the unique solution of the linear equation

$$\Theta'(w^*, \tau_*) \begin{pmatrix} \Delta \hat{w} \\ \Delta \hat{\tau} \end{pmatrix} = -\Theta_{\hat{\sigma}_*}(w^*, \tau_*), \tag{21}$$

cf. (11). Using $\{\hat{\alpha}_k\}_K \to 0$ and taking the limit $k \to \infty$ on the subset K, we then obtain from (18) and (20) that

$$\|\theta(x^*, s^*, \tau_*)\| \ge \beta \tau_* > 0.$$
 (22)

On the other hand, we get from (20), (18), Lemma 3.2 (c), and $\hat{\sigma}_k \leq \hat{\sigma}_{\text{max}}$ that

$$\begin{aligned} \|\theta(\hat{x}^k + \hat{\alpha}_k \Delta \hat{x}^k, \hat{s}^k + \hat{\alpha}_k \Delta \hat{s}^k, (1 - \hat{\sigma}_k \hat{\alpha}_k) \hat{\tau}_k \| &> (1 - \hat{\sigma}_k \hat{\alpha}_k) \beta \hat{\tau}_k \\ &= (1 - \hat{\sigma}_k \hat{\alpha}_k) \beta \tau_k \\ &\geq (1 - \hat{\sigma}_k \hat{\alpha}_k) \|\theta(x^k, s^k, \tau_k)\| \\ &\geq (1 - \hat{\sigma}_{\max} \hat{\alpha}_k) \|\theta(x^k, s^k, \tau_k)\| \end{aligned}$$

for all $k \in \mathbb{N}$ sufficiently large. Using (18) and $\Delta \hat{\tau}_k = -\hat{\sigma}_k \hat{\tau}_k$ (cf. (8)), this implies

$$\frac{\|\theta(x^k + \hat{\alpha}_k \Delta \hat{x}^k, s^k + \hat{\alpha}_k \Delta \hat{s}^k, \tau_k + \hat{\alpha}_k \Delta \hat{\tau}_k)\| - \|\theta(x^k, s^k, \tau_k)\|}{\hat{\alpha}_k} \ge -\hat{\sigma}_{\max} \|\theta(x^k, s^k, \tau_k)\|.$$

Since $\|\theta(\cdot,\cdot,\cdot)\|$ is a continuously differentiable function at (x^*,s^*,τ_*) due to (22), taking the limit $k\to\infty$ for $k\in K$ then gives

$$\frac{\theta(x^*, s^*, \tau_*)^T}{\|\theta(x^*, s^*, \tau_*)\|} \theta'(x^*, s^*, \tau_*) \begin{pmatrix} \Delta \hat{x}^* \\ \Delta \hat{s}^* \\ \Delta \hat{\tau}_* \end{pmatrix} \ge -\hat{\sigma}_{\max} \|\theta(x^*, s^*, \tau_*)\|,$$

where $(\Delta \hat{x}^*, \Delta \hat{\lambda}^*, \Delta \hat{s}^*, \Delta \hat{\tau}_*)$ denotes the solution of the linear system (21). Using (21) then gives

$$-\|\theta(x^*, s^*, \tau_*)\| \ge -\hat{\sigma}_{\max} \|\theta(x^*, s^*, \tau_*)\|.$$

Since $\hat{\sigma}_{\max} \in (0,1)$, this implies $\|\theta(x^*, s^*, \tau_*)\| = 0$, a contradiction to (22).

Note that the assumed existence of an accumulation point in Theorem 3.5 is automatically satisfied under the conditions of Proposition 3.4. — An immediate consequence of Theorem 3.5 is the following result.

Corollary 3.6 Every accumulation point of a sequence $\{w^k\} = \{(x^k, \lambda^k, s^k)\}$ generated by Algorithm 2.1 is a solution of the optimality conditions (3).

Proof. Let $w^* = (x^*, \lambda^*, s^*)$ be an accumulation point of the sequence $\{w^k\} = \{(x^k, \lambda^k, s^k)\}$, and let $\{w^k\}_K$ denote a subsequence converging to w^* . Then we have $\tau_k \to 0$ in view of Theorem 3.5. Hence Lemma 3.2 (c) implies

$$\|\theta(x^*, s^*, 0)\| = \lim_{k \in K} \|\theta(x^k, s^k, \tau_k)\| \le \beta \lim_{k \in K} \tau_k = 0,$$

i.e., we have $x^* \geq 0$, $s^* \geq 0$ and $x_i^* s_i^* = 0$ for i = 1, ..., n due to the definition of θ . Since Lemma 3.2 (a) also shows that we have $A^T \lambda^* + s^* = c$ and $Ax^* = b$, we see that $w^* = (x^*, \lambda^*, s^*)$ is indeed a solution of the optimality conditions (3).

We next want to give a local convergence result. To this end, we first note that the search direction we obtain in our predictor step is identical to the one obtained in the predictor step of the method from [9].

Lemma 3.7 The vector $(\Delta w^k, \Delta \tau_k)$ is a solution of the linear system (10) if and only if Δw^k solves the system

$$\Phi_{\tau_k}'(w^k)\Delta w = -\Phi_0(w^k),$$

and $\Delta \tau_k = 0$.

Proof. Since the smoothing parameter on the right-hand side of the linear system (10) is equal to zero, the assertion follows immediately from the discussion following (6). \Box

The previous result implies that we can apply the local rate of convergence analysis from [9]. Hence we obtain the following result from [9] (see also Tseng [17]).

Theorem 3.8 Assume that the sequence $\{w^k\}$ generated by Algorithm 2.1 converges to a strictly complementary solution of the optimality conditions (2.1). Suppose further that the parameter β from Step (S.0) of Algorithm 2.1 is chosen sufficiently large such that $\beta > 2\sqrt{n}$. Then the predictor step is eventually accepted, and we have

$$\tau_{k+1} = O(\tau_k^2)$$

for all $k \in \mathbb{N}$ sufficiently large, i.e., the smoothing parameter converges locally Q-quadratically to zero.

Note that a typical interior-point method can guarantee the convergence of the corresponding iteration sequence to a strictly complementary solution, so from this point of view, the assumptions we use in Theorem 3.8 are stronger. However, this is basically the only difference, in particular, we stress that the assumptions used in Theorem 3.8 do not necessarily imply that the solution set of the optimality conditions (3) reduces to a singleton.

We close this section by noting that all results (with the possible exception of Theorem 3.8) would still be true if φ would denote the Fischer-Burmeister function

$$\varphi(a,b) := a + b - \sqrt{a^2 + b^2}$$

from [10] together with its smooth counterpart

$$\varphi_{\tau}(a,b) := a + b - \sqrt{a^2 + b^2 + 2\tau^2}$$

from [13]; this can be seen by an easy inspection of the previous proofs. On the other hand, it is currently an open question whether or not Theorem 3.8 also holds for the Fischer-Burmeister function.

4 Numerical Results

We implemented Algorithm 2.1 in MATLAB by modifying the LIPSOL code from Zhang [19, 20]. LIPSOL is a primal-dual interior-point solver for linear programs, written in MATLAB and calling a FORTRAN subroutine in order to solve certain linear systems using the sparse Cholesky method by Ng and Peyton [14]. Since the linear systems occurring in Algorithm 2.1 have essentially the same structure as those arising in primal-dual interior-point methods, it was possible to use the numerical linear algebra part from LIPSOL for our implementation of Algorithm 2.1.

The starting point $w^0 = (x^0, \lambda^0, s^0)$ was constructed in the following way:

- (a) Solve $AA^Ty = b$ using a sparse Cholesky code in order to compute $y^0 \in \mathbb{R}^m$.
- (b) Set $x^0 := A^T y^0$.
- (c) Solve $AA^T\lambda = Ac$ using a sparse Cholesky code to compute $\lambda^0 \in \mathbb{R}^m$.
- (d) Set $s^0 := c A^T \lambda^0$.

This is exactly the starting point used in [9]. One arrives at this starting point by solving the two simple programs

$$\min \frac{1}{2} ||x||^2 \quad \text{s.t.} \quad Ax = b$$

for x^0 and

$$\min \frac{1}{2} ||s||^2 \quad \text{s.t.} \quad A^{\scriptscriptstyle T} \lambda + s = c$$

for λ^0 and s^0 . The construction of the starting point guarantees that the two linear systems Ax = b and $A^T\lambda + s = c$ are satisfied in $w^0 = (x^0, \lambda^0, s^0)$. Furthermore, the initial smoothing parameter τ_0 is taken such that

$$\tau_0 \ge \sqrt{x_i^0 s_i^0} \quad \forall i \in \{1, \dots, n\} \text{ with } x_i^0 > 0, s_i^0 > 0.$$

This choice guarantees that we have $\theta(x^0, s^0, \tau_0) \leq 0$ (both for the minimum and the Fischer-Burmeister function). This condition is required by the algorithm from Burke and Xu [2], although it is not necessary for our method.

We terminate our iteration if one of the following conditions hold:

- (a) $\tau_k < 10^{-4}$ or
- (b) $\|\Phi(w^k)\|_{\infty} < 10^{-4}$ or
- (c) $\|\Phi(w^k)\|_{\infty} < 10^{-3}$ and $\|\Phi(w^k)\|_{\infty}/\|\Phi(w^0)\|_{\infty} < 10^{-6}$.

Criterion (a) was used in [9] and is motivated by the fact that the square of τ does, more or less, play the role of the duality gap in interior-point methods (cf. (4)) for which 10^{-8} is a typical value for the stopping parameter. Criterion (b) is an absolute error measuring the total residual $\|\Phi(w^k)\|_{\infty}$, whereas (c) is a mixture between a weakened form of this absolute error and a relative error comparing the kth residual $\|\Phi(w^k)\|_{\infty}$ with the initial residual $\|\Phi(w^0)\|_{\infty}$.

The remaining parameters from Step (S.0) of Algorithm 2.1 were chosen as follows:

$$\rho = 0.9, \, \beta := \|\Phi_{\tau_0}(w^0)\|/\tau_0$$

and φ being the Fischer-Burmeister function (according to our experience, the Fischer-Burmeister function gives better results than the minimum function, at least within the framework of Algorithm 2.1). Finally, the parameter $\hat{\sigma}_k$ from Step (S.3) of Algorithm 2.1 was always taken to be 0.5.

All test runs were done on a SUN Ultra 2 with 300 MHz, and Table 1 contains the corresponding results, with the columns of Table 1 having the following meanings:

problem: name of the test problem in the netlib collection, number of equality constraints (after preprocessing), m: number of variables (after preprocessing), n: number of iterations until termination, k: P: number of accepted predictor steps, value of τ_k at the final iterate, τ_f : value of $\|\Phi(w^k)\|_{\infty}$ at the final iterate, $\|\Phi(w^f)\|_{\infty}$: primal objective: value of the primal objective function at final iterate.

Table 1: Numerical results for Algorithm 2.1

$\operatorname{problem}$	m	n	k	Р	$ au_f$	$\ \Phi(w^f)\ _{\infty}$	primal objective
25fv47	798	1854	34	17	6.0250e - 04	2.9537e - 04	5.5018459053e+03
80bau3b	2235	11516	29	23	1.0987e - 03	7.1465e - 04	9.8722419211e + 05
adlittle	55	137	15	15	2.1737e - 02	8.7395e - 05	2.2549496391e+05
afiro	27	51	10	10	4.7999e - 02	3.2452e - 05	-4.6474687177e+02
agg	488	615	23	20	$1.3406e{-02}$	6.9052e - 04	-3.5991767286e+07
agg2	516	758	25	18	7.6969e - 03	4.7607e - 04	-2.0239252355e+07

Table 1 (continued): Numerical results for Algorithm 2.1

problem	m	n	k	Р	$ au_f$	$\ \Phi(w^f)\ _{\infty}$	primal objective
agg3	516	758	30	14	7.0617e - 03	1.1522e - 04	1.0312115936e + 07
bandm	269	436	20	19	4.9684e - 04	9.3110e - 05	-1.5862801756e+02
beaconfd	148	270	18	15	2.7426e - 03	5.8563e - 04	3.3592485986e + 04
blend	74	114	13	12	2.7468e - 03	2.9106e - 06	-3.0812134385e+01
bnl1	632	1576	26	16	3.5355e - 04	7.0895e - 05	1.9776295617e + 03
bnl2	2268	4430	26	14	9.5617e - 04	4.8772e - 04	1.8112367543e + 03
boeing1	347	722	26	16	2.2843e - 03	$4.9528e{-04}$	-3.3521310546e+02
boeing2	140	279	16	15	5.4054e - 03	9.8945e - 04	-3.1500732408e+02
bore3d	199	300	28	22	1.3979e - 03	4.1970e - 05	1.3730804026e + 03
brandy	149	259	19	15	1.1040e - 03	7.8205e - 05	1.5185099104e + 03
capri	267	476	20	19	$9.5525e{-03}$	6.4732e - 04	2.6900133856e + 03
cycle	1801	3305	39	19	$9.4566e{-05}$	$1.0106e{-02}$	-5.2249915841e+00
czprob	737	3141	22	19	1.1163e - 02	2.8069e - 04	2.1851966995e+06
d2q06c	2171	5831	57	19	8.3779e - 05	$4.0045e{-05}$	1.2278421095e + 05
d6cube	404	6184	25	21	1.7077e - 03	2.6014e - 05	3.1549167161e + 02
degen2	444	757	23	23	$2.3842e{-03}$	9.9759e - 05	-1.4351779632e+03
degen3	1503	2604	16	16	7.9692e - 04	5.7716e - 05	-9.8729398786e+02
dfl001	6071	12230					
e226	220	469	27	25	2.4792e - 04	5.3902e - 05	-1.8751928739e+01
etamacro	357	692	26	13	1.5436e - 04	7.3792e - 05	-7.5571522983e+02
fffff800	501	1005	36	14	6.2879e - 03	$8.5460e{-04}$	5.5567957590e + 05
finnis	492	1014	31	20	1.3882e - 03	2.9195e - 04	1.7279127031e + 05
fit1d	24	1049	20	18	5.7480e - 04	$4.0534e{-05}$	-9.1463780917e+03
fit1p	627	1677	19	19	1.7472e - 03	7.3692e - 06	9.1463780936e + 03
fit2d	25	10524	22	20	6.0248e - 04	8.2675e - 05	-6.8464293289e+04
fit2p	3000	13525	20	20	1.2942e - 03	3.0570e - 04	6.8464293283e+04
forplan	135	463	28	17	$4.7384e{-03}$	9.3267e - 04	-6.6421820761e+02
ganges	1137	1534	25	20	$3.0644e{-03}$	6.4436e - 04	-1.0958573612e+05
gfrd-pnc	600	1144	23	16	$1.2681e{-02}$	2.4942e - 04	6.9022360024e+06
greenbea	2318	5424	25	20	5.7593e - 03	8.5643e - 04	-7.2462520306e+07
greenbeb	2317	5415	35	15	$2.2551e{-03}$	4.7930e - 04	-4.3022602607e+06
grow15	300	645	37	18	2.4283e - 02	3.7293e - 06	-1.0687094129e+08
grow22	440	946	37	15	$1.0882e{-01}$	7.5577e - 06	-1.6083433646e+08
grow7	140	301	34	19	3.2322e - 02	2.2135e - 05	-4.7787811813e+07
israel	174	316	27	17	3.9926e - 03	2.8436e - 04	-8.9664482178e+05
kb2	43	68	32	10	3.3160e - 03	$9.8866e{-05}$	-1.7499000911e+03
lotfi	151	364	35	16	3.7591e - 03	8.0923e - 04	-2.5263066012e+01
maros	835	1921	37	12	3.1239e - 03	7.1513e - 04	-5.8063742927e+04
maros-r7	3136	9408	22	22	4.7684e - 03	9.7763e - 04	1.4971851671e + 06
modszk1	686	1622	26	17	1.1499e - 02	8.3608e - 04	3.2061981508e + 02
nesm	654	2922	52	14	2.9135e - 04	2.4794e - 04	1.4076036489e + 07
perold	625	1530	33	14	3.3115e - 03	7.3875e - 04	-9.3805322461e+03

Table 1 (continued): Numerical results for Algorithm 2.1

problem	m	n	k	Р	$ au_f$	$\ \Phi(w^f)\ _{\infty}$	primal objective
pilot	1441	4657	81	14	1.1707e - 04	1.4059e - 04	-5.5748445014e+02
pilotja	924	2044	76	14	9.6009e - 05	4.3618e - 03	-6.1130052461e+03
pilotwe	722	2930	61	13	5.8582e - 04	3.7472e - 04	-2.7201075333e+06
pilot4	402	1173	132	13	9.5377e - 05	6.9395e - 03	-2.5810606602e+03
pilot87	2030	6460	63	14	8.3070e - 05	8.8733e - 03	3.0173031374e + 02
$\operatorname{pilotnov}$	951	2242	27	22	2.5409e - 03	3.0714e - 04	-4.4972761773e+03
recipe	85	177	14	14	1.2207e - 03	$3.5042e{-05}$	-2.6661598322e+02
sc105	105	163	19	13	1.4451e - 03	7.7940e - 05	-5.2202033312e+01
sc205	205	317	22	19	7.5991e - 04	1.3473e - 04	-5.2202035425e+01
sc50a	49	77	15	10	3.6860e - 03	4.9576e - 05	-6.4575009902e+01
sc50b	48	76	14	11	6.6556e - 03	7.6186e - 06	-6.9999776566e+01
scagr25	471	671	19	17	$2.3406e{-02}$	2.9238e - 04	-1.4753433056e+07
scagr7	129	185	19	18	3.4159e - 03	3.3656e - 04	-2.3313898243e+06
scfxm1	322	592	20	19	5.9750e - 03	6.4703e - 04	1.8416759818e + 04
scfxm2	644	1184	26	18	4.3503e - 03	8.7736e - 04	3.6660262213e+04
scfxm3	966	1776	26	21	$4.9124e{-03}$	9.6075e - 04	5.4901255716e + 04
scorpion	375	453	21	20	2.7373e - 04	1.8825e - 05	1.8781248227e + 03
scrs8	485	1270	21	19	6.6791e - 04	$4.5185e{-05}$	9.0429695560e + 02
scsd1	77	760	22	22	$4.7684e{-03}$	9.5696e - 06	8.6666991041e+00
scsd6	147	1350	15	15	4.0199e - 04	1.1125e - 06	5.0500000067e+01
scsd8	397	2750	13	13	$1.1068e{-02}$	$4.6656e{-05}$	9.0500023711e+02
sctap1	300	660	24	23	4.0019e - 03	5.1964e - 05	1.4122500207e+03
$\operatorname{sctap} 2$	1090	2500	18	16	1.5629e - 03	1.4780e - 05	1.7248071430e + 03
sctap3	1480	3340	18	17	2.1396e - 03	9.2047e - 05	1.4240000008e+03
seba	515	1036	23	15	3.1158e - 03	9.3931e - 05	1.5711600096e+04
share1b	112	248	43	14	3.0326e - 03	9.3991e - 04	-7.6589318369e+04
share2b	96	162	16	16	3.0518e - 04	$4.4814e{-07}$	-4.1573224024e+02
shell	496	1487	22	14	1.7418e - 01	1.6486e - 05	1.2088253461e+09
ship04l	356	2162	20	20	1.6465e - 02	8.0608e - 04	1.7933245380e + 06
ship04s	268	1414	20	20	1.2018e - 02	$6.5490 \mathrm{e}{-05}$	1.7987147004e + 06
ship08l	688	4339	21	20	1.0490e - 02	3.7678e - 04	1.9090552114e + 06
ship08s	416	2171	20	20	1.8916e - 02	1.4499e - 04	1.9200982105e+06
ship12l	838	5329	21	20	8.4547e - 03	5.6258e - 04	1.4701879193e + 06
ship12s	466	2293	20	19	$8.9471e{-03}$	6.2617e - 04	1.4892361344e + 06
sierra	1222	2715	22	19	1.9638e - 02	9.5043e - 04	1.5394362263e + 07
stair	356	538	19	18	2.3050e - 03	1.6815e - 04	-2.5126689656e+02
standata	359	1258	13	12	7.2079e - 02	$9.5241e{-06}$	1.2577586668e + 03
standgub	361	1366	13	12	7.2079e - 02	$9.5241e{-06}$	1.2577586668e + 03
$\operatorname{standmps}$	467	1258	18	14	9.0258e - 03	1.1631e - 05	1.4060176463e + 03
stocfor1	109	157	16	11	$3.3401e{-02}$	1.2536e - 04	-4.1131976111e+04
stocfor2	2157	3045	29	16	8.3230e - 04	2.4732e - 05	-3.9024408532e+04
stocfor3	16675	23541	63	17	1.8715e - 04	2.8605e - 04	-3.9976784284e+04

Table 1 (continued): Numerical results for Algorithm 2.1

$\operatorname{problem}$	m	n	k	Р	$ au_f$	$\ \Phi(w^f)\ _{\infty}$	primal objective
stocfor3old	16675	23541	70	13	8.7370e - 05	6.0456e - 04	-3.9976783942e+04
truss	1000	8806	19	18	6.3715e - 03	5.6972e - 05	4.5881584778e + 05
tuff	292	617	32	16	3.2629e - 04	$1.5624e{-04}$	$2.9216987102\mathrm{e}{-01}$
vtpbase	194	325	19	19	3.8147e - 02	3.2374e - 05	1.2983146617e + 05
wood1p	244	2595	13	13	3.3617e - 04	$6.1025 \mathrm{e}{-05}$	1.4429024524e+00
woodw	1098	8418	34	22	$2.2390e{-04}$	9.7122e - 05	1.3044869516e + 00

The overall results are quite good and seem to be better than the corresponding results from the three-step method described in [9]. The method has only one failure on problem df1001 (interestingly, LIPSOL also produces an error for this example, at least on our machine), and most test problems can be solved in less than 20–30 iterations. Although interior-point methods are still more efficient on most examples, the numerical behaviour of our smoothing-type method is getting pretty close to the one of interior-point methods, and is definitely approaching an area where it may be viewed as a possible alternative to interior-point methods.

When comparing the results with an interior-point solver, however, one should take into account that Algorithm 2.1 has to factor up to two linear systems of equations per iteration, whereas interior-point methods work with only one factorization. On the other hand, we stress that Algorithm 2.1 has to factorize only one linear system at those iterations where the predictor step is not successful. Moreover, it seems possible to modify the theory in such a way that one can skip the corrector step whenever the predictor step is acceptable. Such a modification of Algorithm 2.1 would then have to factorize only one system per iteration.

5 Concluding Remarks

In this paper, we modified the recently proposed smoothing-type methods from [2, 9]. The modified method has some stronger global and/or local convergence properties than the methods from [2, 9], and the numerical results indicate that the method works very well on the netlib test problem collection. Since these results were obtained by using the Fischer-Burmeister function (rather than the minimum function) and since the local convergence result from Theorem 3.8 does not necessarily hold for the Fischer-Burmeister function, this function certainly deserves further investigation. In fact, this is part of our future research, and we hope that this, in turn, will have a positive influence on our implementation of the predictor step.

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